

Large-deviation properties of the largest biconnected component for random graphs

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Abstract. We study the size of the largest biconnected components in sparse Erdős-Rényi graphs with finite connectivity and Barabási-Albert graphs with non-integer mean degree. Using a statistical-mechanics inspired Monte Carlo approach we obtain numerically the distributions for different sets of parameters over almost their whole support, especially down to the rare-event tails with probabilities far less than 10^{-100} . This enables us to observe a qualitative difference in the behavior of the size of the largest biconnected component and the largest 2-core in the region of very small components, which is unreachable using simple sampling methods. Also, we observe a convergence to a rate function even for small sizes, which is a hint that the large deviation principle holds for these distributions.

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1 Introduction

The robustness of networks [1–5] attracted much interest in recent time, from practical applications for, e.g., power grids [6–8], the internet [9,10], to examinations of genomes [11,12]. As typical in network science, one does not only study the properties of existing networks. To model the properties of real networks, different ensembles of random graphs were devised, e.g., Erdős-Rényi random graphs [13], small world graphs [14], or scale-free graphs [15]. Also for such ensembles the robustness has been studied by analytical and numerical means [16–19]. One often used approach to determine the robustness of networks are *fragmentation* studies, where single nodes are removed from the network. These nodes are selected according to specific rules (“attack”) or randomly (“failure”). The functionality, e.g., whether it is still connected, is tested afterwards. A property necessary for robustness is thus that the graph stays connected when removing an arbitrary node. This exact concept is characterized by the *biconnected component*, which are the connected components which stay connected after an arbitrary node is removed. The existence of a large biconnected component is thus a simple and fundamental property of a graph robust to fragmentation. Another, though related, often studied form of stability looks at the flow through or the transport capability [10] of a graph. Also here a large biconnected component is a good indicator for stability. Intuitively, in a biconnected component there is never a single bottleneck but always a backup path to reach any node. This ensures

the function of the network even in case that an arbitrary edge has too low throughput or an arbitrary node of the biconnected component is damaged.

At the same time the biconnected component is a simple concept enabling to some extent its treatment by analytical means for some graph ensembles. For example, the mean size $\langle S_2 \rangle$ of the biconnected component for a graph with a given degree distribution is known [18]. Also the percolation transition of the biconnected component for scale free and Erdős-Rényi graphs is known to coincide with the percolation transition of the single connected component and its finite size scaling behavior is known [20]. Nevertheless, a full description of any random variable is only obtained if its full probability distribution is known. To our knowledge, concerning the size of the biconnected component this has not been achieved so far for any graph ensemble, neither analytically nor numerically.

For few network observables and some graphs ensembles results concerning the probability distributions have been already obtained so far. For the size of the connected component on Erdős-Rényi random graphs analytical results [21] for the rate function exist, i.e., the behavior of the full distribution for large graph sizes N . Numerically it was shown that this is already for relatively small N a very good approximation [22]. Corresponding numerical results for two-dimensional percolation have been obtained as well [22]. Similarly there are numerical, but no analytical works, scrutinizing the size of the related 2-core over most of its support again for Erdős-Rényi random graphs [23].

Since similar results seem not to be available concerning the biconnected components, and given its importance

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for network robustness, this is an omission that we will start to cure with this study. Here, we numerically obtain the probability density function of the size of the largest biconnected component over a large part of its support, i.e., down to probabilities smaller than 10^{-100} . This enables us also to directly observe large deviation properties, and shows strong hints that the large deviation principle holds [24, 25] for this distribution.

The remainder of this manuscript gives definitions of the graph ensembles and the properties of interest, as well as some known results, in section 2.1 and explains the sampling methods needed to explore the tails of the distributions in section 2.2. The results of our simulations and a discussion will follow in section 3. Section 4 summarizes the results.

2 Models and Methods

2.1 Biconnected Components of Random Graphs

A *graph* $G = (V, E)$ is a tuple of a set of nodes V and edges $E \subset V^{(2)}$. A pair of nodes i, j are called *connected*, if there exists a *path* of edges $\{i, i_1\}, \{i_1, i_2\}, \dots, \{i_{k-1}, i_k\}, \{i_k, j\}$ between them. A *cycle* is a closed path, i.e., the edge $\{i, j\}$ exists and i and j are connected in $G' = (V, E \setminus \{i, j\})$. The *connected components* are the maximal disjoint subgraphs, such that all nodes of each subgraph are connected.

A *biconnected component* (sometimes *bicomponent*) of an undirected graph is a subgraph, such that every node can be reached by two paths, which are distinct in every node except the start- and end node. Thus, if any single node is removed from a biconnected component it will still be a connected component. Therefore clearly, each biconnected component is a connected component. We will also look shortly at *bi-edge-connected components*, which are very similar, but the two paths may share nodes as long as they do not share any edge. Note that a biconnected component is always a bi-edge-connected component, but the reverse is not necessarily true. An example is shown in Fig. 1. In this study we will study only the largest biconnected component G_{bi} of two types of random graphs. Note that, while every biconnected component is also a connected component, the largest biconnected component does not need to be a subgraph of the largest connected component G_{cc} , it may be part of another, smaller, connected component. However, its size $S_2 = |G_{\text{bi}}|$ is always smaller or equal than the size of the largest connected component $S = |G_{\text{cc}}|$. Similarly, the size $S_{2\text{-core}}$ of the largest connected component of the *2-core*, the subgraphs that remain after iterative removal of all nodes with degree less than 2, is an upper bound on S_2 , since the 2-core of a graph consists of bicomponents possibly linked by single edges. In Fig. 1 the largest components of each type are visualized for an example connected graph. In fact for the sizes of the largest of the above introduced subgraphs, the following relation holds.

$$S \geq S_{2\text{-core}} \geq S_{2\text{-edge}} \geq S_2. \quad (1)$$

As we will see below, for the ensemble of Erdős-Rényi random graphs in the percolating phase, the distributions of $S_{2\text{-core}}$, $S_{2\text{-edge}}$, and S_2 are actually very similar to each other. One has indeed to inspect the far tails of the distributions to see differences, which also justifies that we study the large-deviation properties here. For the ensemble of Barabási-Albert graphs we study, the same is true. While the distributions of $S_{2\text{-core}}$ and S_2 look very similar in the main region, a qualitative difference is observable in the tail of small components. The difference is even more pronounced than for the ER case, since the general form of the distribution changes qualitatively to a convex shape for $P(S_{2\text{-core}})$.

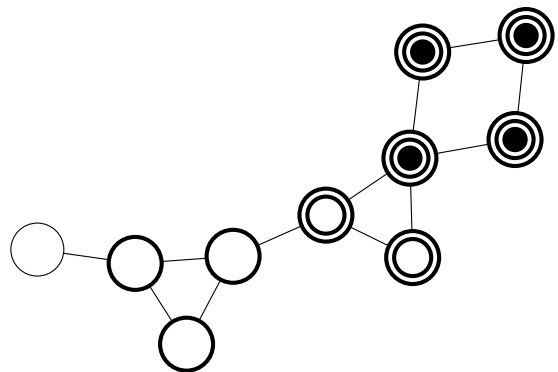


Fig. 1. Every node is part of the connected component, nine nodes with bold borderline are part of the 2-core, six nodes containing a circle are part of the largest bi-edge-connected component and all nodes containing a black dot are part of the largest biconnected component.

The classical way to find biconnected components of a graph [26] is based on a depth first search and thus runs in linear time. For each connected component a depth first search is started at an arbitrary root node of that component. For each node the current *depth* of the search, i.e., at which level in the tree traversed by the depth first search the node is located, and the *lowpoint* are saved. The lowpoint is the minimum of the depth of the neighbors (in the graph) of all descendants of the node (in the tree). If the depth of a node is less or equal the lowpoint of one of its children (in the tree), this node separates two biconnected components and is called *articulation point*. For the root node of the search there is an exception. It is an articulation point, iff it has more than one child. The articulation points separate biconnected components and are members of all biconnected components separated by them. A better illustrated explanation can be found in [27]. After finding all biconnected components, we measure the size of the largest. We did not implement an algorithm, instead we simply used the efficient implementation provided by the LEMON graph library [28].

The mean size of the biconnected component of graphs with a given degree distribution p_k is known for large graphs [18, 29] to be

$$\langle S_2 \rangle = 1 - G_0(u) - (1 - u)G'_0(u), \quad (2)$$

where $G_0(z) = \sum_k p_k z^k$ is the probability generating function, G'_0 its derivative and u the probability to reach a node not part of the giant connected component when following an edge. u is determined by the solution of

$$u = \sum_{k=0}^{\infty} q_k u^k, \quad (3)$$

with the excess degree distribution $q_k = (k+1)p_{k+1}/\langle k \rangle$. Knowing the degree distribution of Erdős-Rényi graphs $G(N, p)$ to be

$$p_k = \binom{N-1}{k} p^k (1-p)^{N-1-k}, \quad (4)$$

allows the numerical evaluation of Eq. (3). We will compare these predictions to our simulational results to scrutinize the behavior for finite N .

The ensemble of *Erdős-Rényi* (ER) graphs $G(N, p)$ consists of N nodes and each of the $N(N-1)/2$ possible edges occurs with probability p . The connectivity $c = Np$ is the average number of incident edges per node, the average *degree*. At $c_c = 1$ this ensemble shows a *percolation transition*. That is in the limit of large graph sizes N the size of the largest connected component is of order $\mathcal{O}(N)$ above this threshold and of order $\mathcal{O}(1)$ below. Interestingly this point is also the percolation transition of the biconnected component [20].

To a lesser extend we also study *Barabási-Albert* (BA) graphs [15]. The ensemble of BA graphs is characterized by a tunable mean degree $\langle k \rangle$ and its degree distribution follows a powerlaw $p(k) \propto k^{-3}$. Realizations are constructed using a growth process. Starting from a fully connected subgraph of m_0 (here $m_0 = 3$) nodes, in every iteration one more node is added and connected to $m \leq m_0$ existing nodes j with a probability $p_j \propto k_j$ dependent on their degree k_j until the size of the graph is N . The parameter $m = 2\langle k \rangle$ by construction. Since $m = 1$ will always result in a tree, which is not biconnected at all and $m \geq 2$ will always be a full biconnected component, we will allow fractional $1 < m < 2$ in the sense, that one edge is always added and a second with probability $m - 1$.

2.2 Sampling

Since we are interested in the far tail behavior of the distribution of the size of the largest biconnected component, it is unfeasible to use naive simple sampling, i.e., uniformly generating configurations, measuring the observable and constructing a histogram. Instead we use a Markov chain Monte Carlo based importance sampling scheme to collect good statistics also in the far tails. This technique was already applied to obtain the distributions over a large range for the score of sequence alignments [30–32], to obtain statistics of the convex hulls of a wide range of types of random walks [33–36], to work distributions for non-equilibrium systems [37] and especially to different properties of Erdős-Rényi random graphs [38, 22, 23, 39].

The Markov chain in this case is a chain of random number vectors ξ_t , $t = 1, 2, \dots$. Each entry of ξ_t is drawn from a uniform $U(0, 1)$ distribution. Each vector serves as an input for a function which generates a random graph. Since all randomness is included in ξ_t , the generated graph $G_t = G(\xi_t)$ depends deterministically on ξ_t . In this way, the Markov chain $\{\xi_t\}$ corresponds to a Markov chain $\{G_t\}$ of graphs. This approach, of separating the randomness from the actually generated objects, has the advantage that for the Markov chain we can generate graph realizations of arbitrary ensembles from scratch, without having to invent a valid Markov chain change move for each ensemble. However, for Erdős-Rényi graphs, we use a specialized change move for performance reasons. One change move is to select a random node i , delete all incident edges and add every edge $\{i, j\}$ with $j \in V \setminus \{i\}$ with probability p . For the Barabási-Albert graphs such a simple change move is not trivial to construct. Therefore, for this type, we perform the typical growth process from scratch after changing one of the underlying random numbers in ξ_t .

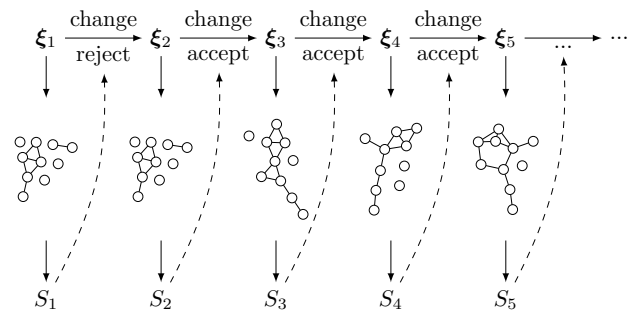


Fig. 2. Four steps of our importance sampling scheme at a small negative temperature, biasing towards a larger biconnected component.

The main idea to obtain good statistics over a large part of the support, especially for probabilities smaller than, say, 10^{-100} , is to bias the generated samples towards those regions. Therefore, we will use classical Metropolis sampling to gather realizations of graphs G . The Markov chain underlying this method consists of either Graph realizations G (ER case) or random number vectors ξ from which a graph realization can be constructed $G(\xi)$ (BA case). We will describe the process for the latter more general case. We start our Markov chain with some random state ξ_1 and at every iteration we propose a new state ξ' , i.e., replace a single entry of ξ_t with a new uniform random number and generate a new realization $G(\xi')$ from these random numbers. We will accept this proposal as the new state ξ_{t+1} , with the classical Metropolis acceptance probability $p_{\text{acc}} = \min\{1, e^{-\Delta S/T}\}$. This process is sketched in Fig. 2. Since we are interested in the size of the largest biconnected component S , we will treat this observable as the “energy” of the realization. Thus, ΔS is the difference in energy between the old and proposed state. Otherwise the proposal is rejected, i.e., $\xi_{t+1} = \xi_t$.

Following this protocol, the Markov chain will equilibrate eventually and from thereon yield realizations $G(\xi)$ which are Boltzmann distributed with respect to some “artificial temperature” T

$$Q_T(G) = \frac{1}{Z_T} e^{-S(G)/T} Q(G), \quad (5)$$

where $Q(G)$ is the natural distribution of the realizations and Z_T the partition function, i.e., a normalization constant. Now, we can use the temperature T as a tuning parameter to adjust the part of the distribution we want to gather samples from. Low positive temperatures will bias the “energy” S towards smaller values because decreases in S are always accepted and increases in S are more often rejected. For negative T this bias works in the opposite way towards larger values of S , i.e., larger biconnected components in this case.

For any chosen temperature, the sampling will be restricted to some interval determined by the value of T . Thus, to obtain the desired distribution $P(S)$ over a large range of the support, simulations for many different temperatures have to be performed. We have to choose the temperatures T in a certain way, to be able to reconstruct the wanted distribution $P(S)$ from this data. First, we can transform $Q(G)$ into $P(S)$ by summing all realizations G , which have the same S . Hence we obtain with Eq. 5

$$\begin{aligned} P_T(S) &= \sum_{\{G|S(G)=S\}} Q_T(G) \\ &= \sum_{\{G|S(G)=S\}} \frac{\exp(-S/T)}{Z_T} Q(G) \\ &= \frac{\exp(-S/T)}{Z_T} P(S). \end{aligned}$$

With this relation we can calculate the wanted, unbiased distribution $P(S)$ from measurements of our biased distributions $P_T(S)$. The ratios of all constants Z_T can be obtained by enforcing continuity of the distribution $P(S)$, i.e.,

$$P_{T_j}(S) e^{S/T_j} Z_{T_j} = P_{T_i}(S) e^{S/T_i} Z_{T_i}.$$

This requires that our measurements for $P_T(S)$ are at least pairwise overlapping such that there is no unsampled region between sampled regions. From pairwise overlaps the pairwise ratios Z_i/Z_j can be approximated. The absolute value of the Z_T can afterwards be obtained by the normalization of $P(S)$. Although the size of the largest biconnected component S_2 is a discrete variable for every finite N and should therefore be normalized such that the probabilities of every event should sum to one, i.e., $\sum_{i=0}^N p(S_2 = i/N) = 1$, we are mainly interested in the large N behavior and especially the rate function. This limit is continuous and should therefore be treated with a different normalization $\int_0^1 p(S_2) dS_2 = 1$, which we approximate for finite N by the trapezoidal rule. Anyway, the difference here is just a factor N .

While this technique does usually work quite well and all distributions but one exception are obtained with this

method, there are sometimes first order phase transitions within the finite temperature ensemble, rendering it infeasible, or at least very tedious, to acquire values inbetween two temperatures. This was a problem here for the modified Barabási-Albert graph at the largest simulated graph size N . This phenomenon is well known and explored in detail in [22]. We filled this gap by modified Wang-Landau simulations [40–44] with subsequent entropic sampling [45, 46].

3 Results

We applied the temperature-based sampling scheme to ER with finite connectivities of $c \in \{0.5, 1, 2\}$ and BA with $m = 1.3$ over practically the whole support $S_2 \in [0, 1]$ using around a dozen different temperatures for each ensemble and Markov chains of length $10^6 N$ to gather enough samples after equilibration and discarding correlated samples. Additionally for BA the range $S_2 \in [0.1, 0.35]$ was sampled using Wang-Landau’s method and merged into results obtained from the temperature based sampling for the remainder of the distribution. All error estimates for the distributions are obtained via bootstrap resampling [47, 48] but are always smaller than the symbol size and therefore not shown. Error estimates for fits are Gnuplot’s asymptotic standard errors.

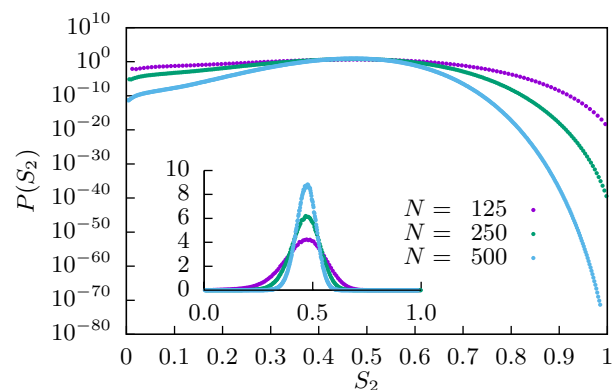


Fig. 3. Distributions of the size of the largest biconnected component S_2 for ER graphs at connectivity $c = 2$ and three different graph sizes N . The main plot shows the distributions in logarithmic scale to display the tails, the inset shows the same distributions in linear scale, where a concentration around the mean value $\langle S_2 \rangle$ (cf. Fig 4) is visible. (For clarity not every bin is visualized.)

Examples for the distribution of the largest biconnected component’s size for ER graphs at $c = 2.0$ are depicted in Fig. 3 at three different graph sizes N . The inset shows the distribution in linear scale, where a concentration with increasing size N around the mean value is visible. While the main part of the distribution in the inset looks rather symmetric, the tails are obviously not. Also it is visible that the tails of the distribution get more and more suppressed when increasing the value of N .

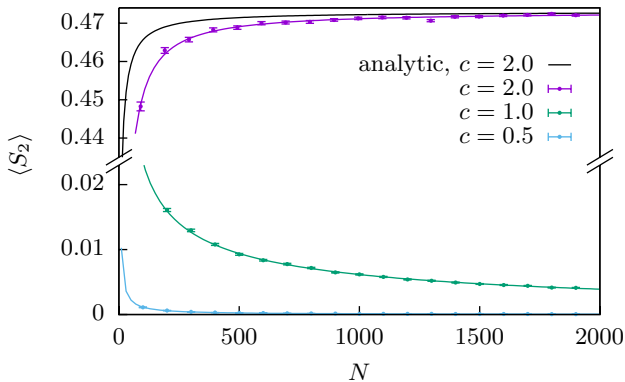


Fig. 4. Mean size of the largest biconnected component $\langle S_2 \rangle$ for different graph sizes N . Notice the broken $\langle S_2 \rangle$ -axis. The black line denotes the analytic expectation for $c = 2$ from Eq. 2 [18]. The expectation for $c \leq 1$ is $\langle S_2 \rangle = 0$. Fits to a power law with offset $\langle S_2 \rangle = aN^b + S_2^\infty$ lead to $S_2^\infty = -6(8) \cdot 10^{-6}$ for $c = 0.5$, $S_2^\infty = -0.0013(4)$ for $c = 1.0$ and $S_2^\infty = 0.4729(3)$ for $c = 2.0$.

Since the mean size of the biconnected component of ER is known for large enough graphs, we will compare the mean sizes of our simulations to the analytical expectation. Those results are shown in Fig. 4, notice the broken $\langle S_2 \rangle$ -axis. Apparently at $c = 2$ for small sizes N the analytical approximation, while close to our measurements, overestimates the size of the biconnected component slightly but the relative error diminishes for larger sizes. In fact, we extrapolated our measurements to the limit of large N using a power-law ansatz $\langle S_2 \rangle = aN^b + S_2^\infty$ yielding for $c = 0.5$ an offset S_2^∞ compatible within errorbars with the expectation $\langle S_2 \rangle = 0$ (exact values in the caption of Fig. 4), which is quite remarkable for our ad-hoc fit function. The case $c = 1$ suggests a negative S_2^∞ close to zero, which is probably caused by correction to our assumed scaling law. The case $c = 2$ seems to converge to the limit of the analytical expectation also.

To compare the sizes of different relevant types of components, Fig. 5 shows the distributions of the relative size of the largest connected component S [22], the largest 2-core $S_{2\text{-core}}$ [23], the largest bi-edge-connected component $S_{2\text{-edge}}$ and the largest biconnected component S_2 for $N = 500$ and $c = 2$ ER graphs. Interestingly, the distributions $P(S_2)$, $P(S_{2\text{-edge}})$ and $P(S_{2\text{-core}})$ are almost identical and only deviate in the region of very small components from each other. As would be expected by the order of Eq. 1, the probability to find very small 2-cores is lower than to find bi-edge-connected components of the same small size, which are again slightly less probable than biconnected components of that size. Anyway, when considering ER graphs, which exhibit by construction no particular structure, the robustness properties which are determined by the biconnected component, can be with very high probability readily inferred from the 2-core.

To understand the topology of the instances of very low probabilities better, we will look at the correlations of the size of the largest connected component S and the largest biconnected component in Fig. 6. Note that this histogram

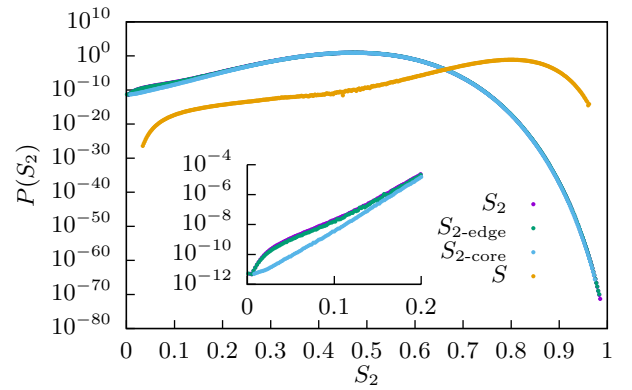


Fig. 5. Comparison of the relative size of the largest connected component S [22], the largest 2-core $S_{2\text{-core}}$ [23], the largest bi-edge-connected component $S_{2\text{-edge}}$ and the largest biconnected component S_2 for $N = 500$ and $c = 2$ ER graphs. The last three are nearly identical for sizes $S_x \gtrsim 0.2$. The inset shows a zoom to the very small components, which is the only region, where the three last types deviate considerably from one another. For clarity not every data point is visualized.

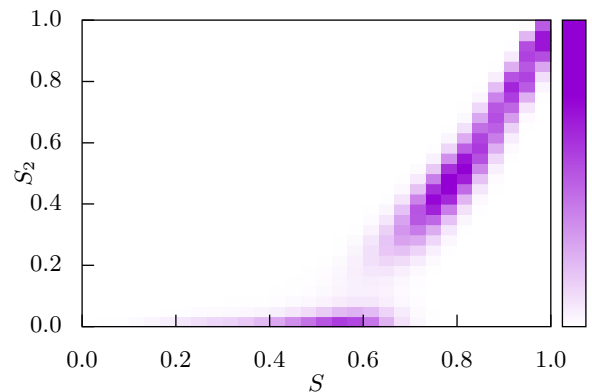


Fig. 6. Correlation histogram of our raw and biased simulation data. A large biconnected component does most probably appear in graphs whose connected component is larger than $S \gtrsim 0.6$.

does not reflect the probabilities, but does count the instances we generated within one of our simulations, i.e., data for many different temperatures are shown without correction for the introduced bias. Anyway, it is instructive to look at this sketch for qualitative understanding. This data is for $c = 2$ ER at $N = 500$. We observe that, even for our biased sampling, there are basically no large biconnected components if the connected component is smaller than $S \lesssim 0.6$. Above this point, we observe larger biconnected components, but generally very few around the size $S_2 \approx 0.2$. Above $S \gtrsim 0.6$ the size of the largest biconnected component is strongly correlated with the size of the largest connected component.

For a qualitative understanding of this behavior, consider the following heuristic argument. For the instances without or with very small biconnected components, i.e., only short cycles, the graph is basically tree-like. Larger biconnected components are then created by connecting

two nodes of the tree with each other, leading to a cycle which is on average roughly half in the order of the size of the tree, leading to the jump in the size of biconnected components. The configurations with smaller biconnected components are apparently entropically suppressed.

Next we will look at the empirical large deviation *rate function* of the measured distributions. The rate function Φ describes the behavior of distributions, whose probability density decays exponentially in the tails in respect to some parameter N . In this case, the parameter N is the graph size. For increasing graph size N the biconnected components which are not typical will be exponentially suppressed. To be more precise, the definition of the rate function $\Phi(S_2)$ is via $P(S_2) = e^{-N\Phi(S_2)+o(N)}$ for the large N limit with the Landau symbol o for terms of order less than N . Since we obtained the distributions over most of their support, we can access the empirical rate functions $\Phi_N(S_2) = -1/N \log P(S_2)$ for finite N . Note that the empirical rate functions do contain all information of the distribution. If for increasing N they converge to the limiting rate function, one says the distribution follows the *large deviation principle* [24, 25].

In Fig. 7 the empirical rate functions for ER at different connectivities c and for different sizes N are shown. Already these comparatively small values of N show remarkably similar empiric rate functions and strongly hint at a convergence to a limit form. While the rate function in the right tail for larger than typical components S_2 are already almost indistinguishable, the convergence seems a bit slower in the left tail of smaller than typical components. This behavior is very similar to the behavior of the sizes of the connected component [22] and the 2-core [23]. This means that, the large deviation principle seems to hold for this distribution.

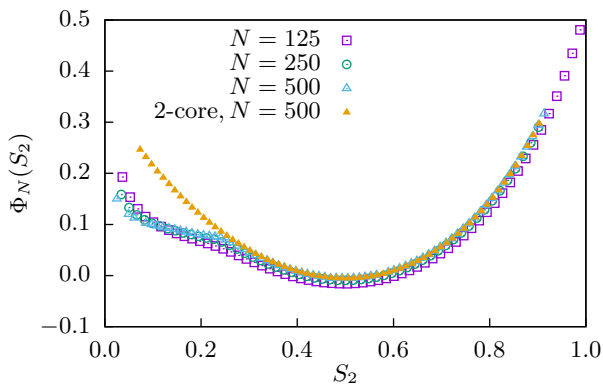


Fig. 8. Empirical rate function $\Phi_N(S_2)$ for multiple graph sizes N of the BA graph ensemble with $m = 1.3$.

The rate function of the largest biconnected component of the BA ensemble at $m = 1.3$ is shown in Fig. 8. The rate function and therefore the distribution does look qualitatively similar to the $c = 2$ case of the ER ensemble (cf. Fig. 7(c)). The dip around $S_2 \approx 0.2$ is more pronounced leading to a more severe discontinuity in the simulated finite temperature ensemble necessitating the use of

Wang-Landau sampling. The main qualitative difference of the behavior of the two distributions is the behavior for very small sizes of the largest biconnected components S_2 , where the empirical rate functions cross each other, hinting at some kind of finite size effect suppressing very small biconnected components in small graphs.

In comparison with the rate function of the 2-core, also shown in Fig. 8, their difference in the region of very small biconnected components respectively 2-cores, which was already observable in ER, is very strong in the BA ensemble. Despite those two observables are almost indistinguishable in the main region, they show strongly different behavior in their overall shape, i.e., the distribution of the 2-core seems convex over the region we obtained statistics for.

4 Conclusions

The biconnected component is the fundamental graph-theoretical concept which is most related to robustness properties of random networks. Nevertheless, the distribution of its size has not been studied before, to our knowledge. We used sophisticated sampling methods to obtain the distributions of the size of the largest biconnected component S_2 , for multiple ER graph ensembles and a modified BA graph ensemble, over a large part of their support.

For the ER ensemble, looking into the large deviation tails of this distribution shows qualitative differences between the size of the 2-core and the biconnected component, which are otherwise not well observable. This is even more the case for the BA ensemble where the overall shape of the distribution seems to differ. While the 2-core distribution seems convex, the distribution of the biconnected component shows a “shoulder”. These qualitative difference, however is only apparent below probabilities of 10^{-20} and are therefore unobservable using conventional methods.

Further, the empirical rate functions are already for the small sizes that we simulated very close to each other hinting at a very fast convergence to the limiting form. Thus, our results indicate that the large deviation principle holds for the numerically obtained distributions. This “well-behaving” of our numerical results may make it promising to address the distribution of the biconnected component by analytical means, which has not been done so far to our knowledge. Furthermore, it would be interesting to study other network ensembles, which are even more relevant for modeling robustness properties, e.g., two-dimensional networks modeling power grids [8] and other transportation networks.

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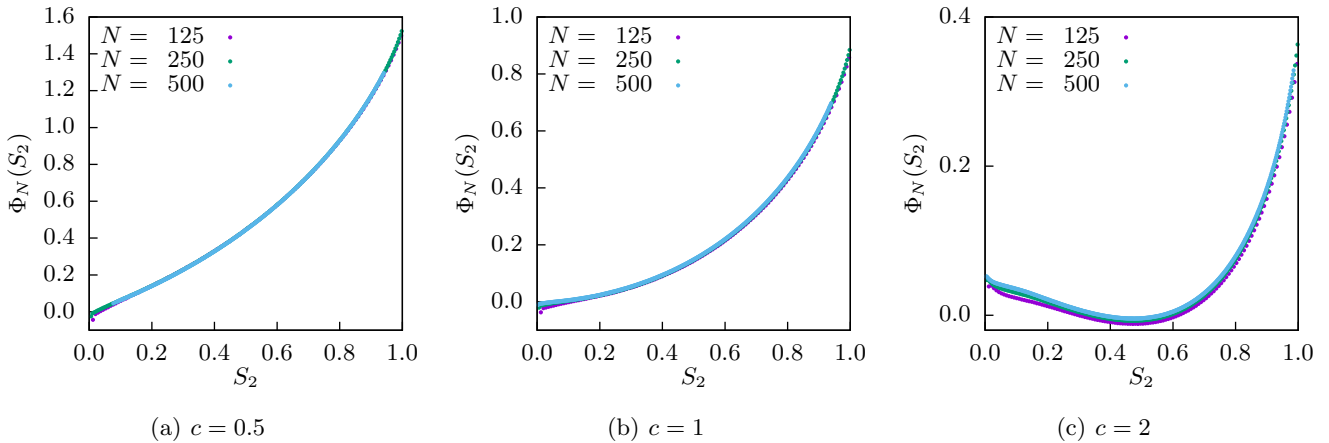


Fig. 7. Empirical rate function $\Phi_N(S_2)$ for multiple graph sizes N and connectivities c of the ER graph ensemble.

6 Authors contributions

AKH conceived the study, HS wrote the first draft of the manuscript and generated most of the new data. All authors contributed ideas, simulation data and analysis to this study. All authors were involved in the preparation of the manuscript.

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